

Convex Optimization

(EE227A: UC Berkeley)

Lecture 8
Weak duality
14 Feb, 2013

○

Suvrit Sra

Primal problem

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$$\begin{aligned} \min \quad & f_0(x) \\ \text{s.t.} \quad & f_i(x) \leq 0, \quad 1 \leq i \leq m, \\ & x \in \{\text{dom } f_0 \cap \text{dom } f_1 \cdots \cap \text{dom } f_m\}. \end{aligned} \tag{P}$$

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- ▶ In our initial derivation: no restriction to convexity.

Lagrangian

To the primal problem, associate **Lagrangian** $\mathcal{L} : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$,

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- ♠ Suppose x is feasible, and $\lambda \geq 0$. Then, we get the lower-bound:

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- ♠ Lagrangian helps write problem in **unconstrained form**

Lagrangian

Claim: Since, $f_0(x) \geq \mathcal{L}(x, \lambda) \quad \forall x \in \mathcal{X}, \quad \lambda \in \mathbb{R}_+^m$, primal optimal

$$p^* = \inf_{x \in \mathcal{X}} \sup_{\lambda \geq 0} \mathcal{L}(x, \lambda).$$

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Proof:

- ♠ If x is not feasible, then some $f_i(x) > 0$
- ♠ In this case, inner sup is $+\infty$, so claim true by definition
- ♠ If x is feasible, each $f_i(x) \leq 0$, so $\sup_{\lambda} \sum_i \lambda_i f_i(x) = 0$

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- ▶ $\forall x \in \mathcal{X}, \quad f_0(x) \geq \inf_{x'} \mathcal{L}(x', \lambda) = g(\lambda)$
- ▶ Now minimize over x on lhs, to obtain

$$\forall \lambda \in \mathbb{R}_+^m \quad p^* \geq g(\lambda).$$

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- ▶ **dual feasible:** if $\lambda \geq 0$ and $g(\lambda) > -\infty$
- ▶ **dual optimal:** λ^* if sup is achieved
- ▶ Lagrange dual is always concave, regardless of original

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Def. Denote **dual optimal value** by d^* , i.e.,

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Theorem (Weak-duality): For problem (P), we have $p^* \geq d^*$.

Proof: We showed that for all $\lambda \in \mathbb{R}_+^m$, $p^* \geq g(\lambda)$.

Thus, it follows that $p^* \geq \sup g(\lambda) = d^*$.

Equality constraints

$$\begin{aligned} \min \quad & f_0(x) \\ \text{s.t.} \quad & f_i(x) \leq 0, \quad i = 1, \dots, m, \\ & h_i(x) = 0, \quad i = 1, \dots, p. \end{aligned}$$

Exercise: Show that we get the Lagrangian dual

$$g : \mathbb{R}_+^m \times \mathbb{R}^p : (\lambda, \nu) \mapsto \inf_x \mathcal{L}(x, \lambda, \nu),$$

where the Lagrange variable ν corresponding to the equality constraints is unconstrained.

Hint: Represent $h_i(x) = 0$ as $h_i(x) \leq 0$ and $-h_i(x) \leq 0$.

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Hint: Represent $h_i(x) = 0$ as $h_i(x) \leq 0$ and $-h_i(x) \leq 0$.

Again, we see that $p^* \geq \sup_{\lambda \geq 0, \nu} g(\lambda, \nu) = d^*$

Some duals

- ▶ Least-norm solution of linear equations: $\min x^T x$ s.t. $Ax = b$
- ▶ Linear programming standard form
- ▶ Study example (5.7) in BV (binary QP)